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A computer-assisted proof of dynamo growth in the stretch-fold-shear map

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ABSTRACT
The Stretch-Fold-Shear (SFS) operator $S_\alpha$ is a functional linear operator acting on complex-valued functions of a real variable $x$ on some domain containing $[-1, 1]$ in $\mathbb{R}$. It arises from a stylized model in kinematic dynamo theory where magnetic field growth corresponds to an eigenvalue of modulus greater than 1. When the shear parameter $\alpha$ is zero, the spectrum of $S_\alpha$ can be determined exactly, and the eigenfunctions corresponding to non-zero eigenvalues are related to the Bernoulli polynomials. The spectrum for $\alpha > 0$ has not been rigorously determined although the spectrum has been approximated numerically. In this paper, a computer-assisted proof is presented to provide rigorous bounds on the leading eigenvalue for $\alpha \in [0, 5]$, showing inter alia that $S_\alpha$ has an eigenvalue of modulus greater than 1 for all $\alpha$ satisfying $\pi/2 < \alpha \leq 5$, thereby partially confirming an outstanding conjecture on the SFS operator.

1. Introduction
In kinematic dynamo theory, the induction equation for the growth of a magnetic field transported by a fluid flow is:

$$\partial_t B = \nabla \times (u \times B) + \varepsilon \nabla^2 B$$

[11, 10, 11]. Here $B$ is the magnetic field, $u$ the fluid flow, $\varepsilon$ the magnetic diffusivity and $\varepsilon^{-1}$ the magnetic Reynolds number. The fluid flow $u$ is assumed to be given, and could be steady or unsteady; it is usually taken to be incompressible, with $\nabla \cdot u = 0$. For example, a seminal study of Arnold and Korkina [4] considers magnetic field evolution in the steady, three-dimensional ABC flow

$$u = (\cos y + \sin z, \cos z + \sin x, \cos x + \sin y).$$

For any such given $u$ as well as a general initial condition, the magnetic growth rate is

$$\gamma(\varepsilon) = \lim_{t \to \infty} t^{-1} \log \|B(t)\|$$

and the flow $u$ is a dynamo if $\gamma(\varepsilon) > 0$ for some diffusivity $\varepsilon > 0$. For example, the above ABC flow is a dynamo as growth occurs for $8.9 \leq \varepsilon^{-1} \leq 10$. 
17.5 and for $\varepsilon^{-1} \geq 27$ up to the limits of present numerical simulations [15]. In astrophysical contexts, the magnetic Reynolds number $\varepsilon^{-1}$, a dimensionless measure of the electrical conductivity, is very large, and this motivates the study of kinematic dynamos in the delicate limit of zero diffusion $\varepsilon \to 0$. If the dynamo growth rate $\gamma(\varepsilon)$ remains positive and of order unity as $\varepsilon$ tends to zero, the flow $\mathbf{u}$ is called a fast dynamo. In this case, we can define the fast dynamo growth rate $\gamma_0$ by $\gamma_0 = \liminf_{\varepsilon \to 0} \gamma(\varepsilon)$.

Delicate aspects of this limit become apparent if we look at the perfectly conducting situation $\varepsilon = 0$, when magnetic field evolution becomes intimately related to the streamline structure of the flow, as may be seen by using $\nabla \cdot \mathbf{u} = \nabla \cdot \mathbf{B} = 0$ to write (1) as

$$\partial_t \mathbf{B} + \mathbf{u} \cdot \nabla \mathbf{B} - \mathbf{B} \cdot \nabla \mathbf{u} \equiv \partial_t \mathbf{B} + \mathcal{L}_\mathbf{u} \mathbf{B} = 0 \quad (\varepsilon = 0). \quad (3)$$

Here the effect of the flow $\mathbf{u}$ on the magnetic field $\mathbf{B}$ is seen to be Lie-dragging: in other words, magnetic field vectors are advected, rotated and stretched with infinitesimal line elements in the flow. Another viewpoint is that magnetic field lines are materially transported in the flow, this being known as Alfvén’s theorem. For a smooth flow $\mathbf{u}$ of any complexity, for example, one that is chaotic in the sense of having a positive topological entropy and so guaranteed exponential growth of line length in the flow, the magnetic energy increases exponentially, $\gamma(0) > 0$, in this case when the diffusion is strictly zero [21].

While it is easy enough to gain magnetic field growth in a chaotic flow with $\varepsilon = 0$, the growing field becomes indefinitely fine-scaled with frequent changes of sign, and the result is that the introduction of any magnetic diffusion $\varepsilon > 0$, no matter how small, can have a dramatic effect on the growth rates of field. For example, a chaotic flow in the periodic plane will have $\gamma(0) > 0$ but it is known through the Zeldovich anti-dynamo theorem [35] that $\gamma(\varepsilon) < 0$ for all $\varepsilon$ and so $\gamma_0 \leq 0$. The singular nature of the $\varepsilon \to 0$ limit versus $\varepsilon = 0$ arises because the stretching is accompanied by the folding of field, and when these folds reach a diffusive scale, the effect of diffusion, no matter how weak, is dramatic and destructive.

Thus flows in the plane cannot be fast dynamos, or dynamos at all, and some three-dimensional motions are needed. Perhaps the simplest class of fast dynamos, based largely on numerical evidence, consists of flows taking the convenient form $\mathbf{u}(x, y, t) = (u_x, u_y, u_z)$, independent of $z$. Field amplification operates via the so-called stretch-fold-shear (SFS) mechanism [5], where magnetic field lines are stretched and folded through the action of $u_x$ and $u_y$, and then shear in the $z$ direction partially aligns the magnetic field vectors so that growth can occur. An example is the Galloway–Proctor ‘circularly polarised’ flow [16],

$$\mathbf{u} = \sqrt{3/2} (\cos(y + \sin t), \sin(x + \cos t), \cos(x + \cos t) + \sin(y + \sin t)). \quad (4)$$

Even in such flows, that in simulations show robust fast dynamo action, the limiting growth rate $\gamma_0$ is found numerically to be strictly less than $\gamma(0)$; in other words, even very weak diffusion has a significant destructive effect on the highly convoluted magnetic fields that are generated, highlighting the delicate aspects of the $\varepsilon \to 0$ limit. We remark that there is only one rigorous result of any generality, of Klapper and Young [21], that the topological entropy of a smooth flow field is an upper bound on the fast dynamo growth rate $\gamma_0$, but there is no rigorous theory on lower bounds for smooth flows in (say) periodic three-dimensional space.
For this reason, a number of stylized models have been considered in fast dynamo theory. Here we consider one designed to illustrate the SFS mechanism, due originally to Childress and Bayly [5] and studied extensively in [17–19], which is

\[ M : [-1, 1]^2 \times \mathbb{R} \to [-1, 1]^2 \times \mathbb{R}, \quad M = M_{Sh} M_{SF}, \]  

(5)

where \( M_{SF} \) is a folded baker’s map:

\[ M_{SF}(x, y, z) = \begin{cases} 
\left( \frac{1}{2} (x - 1), 1 + 2y, z \right), & y < 0, \\
\left( \frac{1}{2} (1 - x), 1 - 2y, z \right), & y \geq 0, 
\end{cases} \]  

(6)

and \( M_{Sh} \) is the shear in the \( z \)-direction: \( M_{Sh}(x, y, z) = (x, y, z + \alpha x) \), where \( \alpha \geq 0 \) is the shear parameter. Here a map is deployed as an idealized model of the motion of points and magnetic field vectors over a single period of a periodic flow such as (4).

Taking \( B(x, y, z) = e^{ix} b(x) \hat{y} + \text{complex conjugate} \), where \( b(x) \) is a complex-valued function of a real variable \( x \) on \([-1, 1]\), the maps \( M_{SF} \) and \( M_{Sh} \) induce maps of \( b \) by the operators:

\[ T_{SF} b(x) = \begin{cases} 
2b(1 + 2x), & x < 0, \\
-2b(1 - 2x), & x \geq 0, 
\end{cases} \quad \text{and} \quad T_{Sh} b(x) = e^{-i\alpha x} b(x). \]  

(7)

This leads to the SFS operator \( T = T_{Sh} T_{SF} \) on \( b(x) \):

\[ T b(x) = \begin{cases} 
2e^{-i\alpha x} b(1 + 2x), & x < 0, \\
-2e^{-i\alpha x} b(1 - 2x), & x \geq 0. 
\end{cases} \]  

(8)

The phase shifts from the shear in Equation (8) generate an ever more complicated field \( T^n b \) as \( n \) increases. To check whether this is a model for a fast dynamo, strictly speaking, some weak diffusive process should be introduced and growth rates studied in the limit \( \varepsilon \to 0 \), of weak diffusion. However, an alternative route is to replace the averaging effect of diffusive smoothing by integration of \( T^n b \) against a smooth test function \( c(x) \), namely to look at the growth of

\[ (c, T^n b) \]  

with the inner product

\[ (f, g) \equiv \int_{-1}^{1} f(x) g(x) \, dx. \]  

(9)

If \((c, T^n b)\) grows exponentially, with growth rate \( \Gamma > 0 \) for general functions \( b \) and \( c \), the map \( T \) is called a perfect dynamo. It is conjectured, with fair numerical evidence, that \( \Gamma = \gamma_0 \); see [10] and references therein. In this way, via this so-called flux conjecture, one aims to link the growth for the fast dynamo limit of weak, non-zero diffusion, with the ever-evolving field for zero diffusion by making use of an appropriate measure of growth, namely integration against a smooth test function. With this in mind, the present paper will focus only on perfect dynamo action in the SFS mapping [17].

Once we look for exponential growth in \((c, T^n b)\), it makes sense to introduce the adjoint operator \( T^* \) (with respect to the above inner product) such that \((c, T^n b) = (T^* c, b)\) for all
This adjoint operator $T^*$ is defined by
\[ T^*c(x) = e^{i\alpha(x-1)/2}c((x-1)/2) - e^{i\alpha(1-x)/2}c((1-x)/2). \] (10)

Note that $T^*$ is the composition of two operators $T^*_{SF}$ and $T^*_{Sh}$, where
\[ T^*_{SF}c(x) = c((x-1)/2) - c((1-x)/2) \quad \text{and} \quad T^*_{Sh}c(x) = e^{i\alpha x}c(x). \] (11)

The key advantage in passing from consideration of the growth of $T^n b$ to growth in $T^*nc$ is that while iterating $T$ on an initial condition creates ever finer oscillations in a complex field $T^n b$, applying $T^*$ repeatedly giving $T^*nc$ stretches out and smooths structure, this being evident from numerical simulations [5].

Thus for the SFS model to be a perfect dynamo (and so by the flux conjecture a fast dynamo) we seek an eigenvalue of $T$ and $T^*$ of modulus greater than unity; moreover, the smoothness of the eigenfunctions of the operator $T$ or $T^*$ is the primary consideration.

Since the operators $T$ and $T^*$ are not compact in $L^2$ and their eigenfunctions are not generally smooth, it is convenient to restrict to a smaller function space which is a subset of $L^2$.

Because of its analyticity-improving property, it is most convenient to study the operator $T^*$ (rather than $T$) restricted to a smaller space $\mathcal{F}$ for which the functions are not only smooth but also analytic in a disc on the complex plane. The restricted operator is denoted by $S$. This restricted SFS operator $S$ is compact and growing smooth eigenfunctions of this operator can be obtained to establish perfect dynamo action in the SFS model. For further details on the SFS dynamo mechanism see [17] and for a discussion of passive scalar decay in baker’s maps, where the effects of diffusion and boundary conditions are more delicate, see [19].

The above considerations lead to the following definition of the SFS operator which is the object of study in this paper.

**Definition 1.1 (The SFS Operator):** The SFS operator $S$ or $S_\alpha$ is defined by
\[ Sc(x) = e^{i\alpha(x-1)/2}c((x-1)/2) - e^{i\alpha(1-x)/2}c((1-x)/2), \] (12)

where $\alpha \geq 0$ is the shear parameter. For $\alpha=0$, $S$ takes on the simpler form
\[ Sc(x) = c((x-1)/2) - c((1-x)/2). \]

Following the above discussion, we seek magnetic flux growth in the form of a perfect dynamo, which corresponds to the existence of an eigenvalue $\lambda$ of $S$, with $|\lambda| > 1$. In [17], it is conjectured that for all $\alpha > \frac{\pi}{2}$, $S$ has at least one eigenvalue $\lambda$ with $|\lambda| > 1$. In fact, numerical evidence suggests that for $\alpha > \frac{\pi}{2}$ the principal eigenvalue satisfies $|\lambda| > 1$ and that, for all families of eigenvalues $\lambda$, $\lim_{\alpha \to \infty} |\lambda| = \sqrt{2}$. Figure 1 (reproduced from the paper [17]) shows the absolute values $|\lambda|$ against $\alpha$.

The purpose of this paper is to study the spectrum of $S$ and, specifically, to prove the following theorem:

**Theorem 1.1 (Computer assisted):** Let $\alpha \in [0,5]$. Then there exists an eigenvalue-eigenfunction pair $(\lambda, c(x))$ of the operator $S$ acting on a Banach space $\mathcal{F}_r$ of analytic functions on the disc $D(0, r)$ of radius $r$ centred on $0$ in $\mathbb{C}$ satisfying the following:
Figure 1. Graph of $|\lambda|$ for eigenvalues $\lambda$ of the SFS operator $S$ with $\alpha$ on the horizontal axis and the modulus of the eigenvalues of $S$ on the vertical axis.

(1) $c(x) \in F_r$, for some $r > 1$.
(2) (a) For $\alpha \in [0, \pi/2)$, $|\lambda| < 1$.
    (b) For $\alpha \in (\pi/2, 5]$, $|\lambda| > 1$.
    (c) For $\alpha = \pi/2$, $\lambda = -i$.

Note that in 2(a) it is not claimed that all eigenvalues have modulus less than one for $\alpha \in [0, \pi/2)$ (although numerical results suggest that this is the case), but rather that there is at least one eigenvalue with this property.

Theorem 1.1 provides a partial proof of a conjecture contained in [17, 18]. Its proof uses computer-assisted estimates to obtain mathematically rigorous results. In the context of kinematic dynamo theory, the theory demonstrates that the SFS model supports the growth of a perfect dynamo for the shear parameter $\alpha > \pi/2$.

2. Properties of the SFS operator

The Stretch-Fold-Shear (SFS) operator is a family $S_\alpha$ of linear operators, depending on the real non-negative parameter $\alpha$, which acts on complex-valued functions $c(x)$ defined on a domain containing the real closed interval $[-1, 1]$. In general, we shall write $S$ for $S_\alpha$ whenever the value of $\alpha$ is clear from the context or of no particular relevance.

For the analysis below, we restrict to $S$ acting on a space $F_r$ of complex-valued functions analytic on the disc $D_r$ about 0 of radius $r$. Here $r > 1$ and may depend on the specific value of the parameter $\alpha$ (for further details see Section 3) but, in any case, $[-1, 1] \subset D_r$. The operator $S$ may be considered on other classes of functions but analytic functions are particularly suited for our purposes.

Because historically the focus is on the real closed interval $[-1, 1]$, we denote the function argument by $x$ rather than $z$. Moreover, depending on the context, we refer to the
function as $c$ or $c(x)$, not restricting $c(x)$ to denote the value of the function at the point $x$. The function spaces $\mathcal{F}_r$ in which we work will be discussed in detail in Section 3.

We note that this operator is well defined for functions on $D_r$, since for all $x \in D_r$, both $(x - 1)/2$ and $(1 - x)/2 \in D_r$ provided $r > 1$. Below we review some properties of this operator, but here we note, in passing, that on the Banach space $\mathcal{F}_r$, $S$ is a compact operator and therefore its spectrum consists of discrete eigenvalues together with 0. We refer to an eigenvector $c(x)$ corresponding to an eigenvalue $\lambda$ of $S$ as an eigenfunction and the pair $(\lambda, c(x))$ as an eigenvalue–eigenfunction pair.

The operator $S$ may readily be rewritten

$$Sc(x) = \left[ e^{i\alpha x} c(x) \right]_{2\alpha} \left( \frac{x - 1}{2} \right),$$

where $[f(x)]_{2\alpha}$ is the function $f(x) - f(-x)$, i.e. twice the odd part of $f$. The form (13) illustrates that for any even function $d_c(x)$ the function $c(x) = e^{-i\alpha x}d_c(x)$ is mapped to 0 by $S$, so that 0 is an eigenvalue of infinite multiplicity. Indeed, functions of this type form the 0-eigenspace of $S$.

Now consider an eigenvalue–eigenfunction pair $(\lambda, c(x))$, where $\lambda \in \mathbb{C}$, $c(x) \in \mathcal{F}_r$ satisfy

$$S(c)(x) = e^{i\alpha(x - 1)/2}c((x - 1)/2) - e^{i\alpha(1-x)/2}c((1 - x)/2) = \lambda c(x).$$

Below, when required, we shall choose a normalization $c(0) = 1$, although often formulae are simplified by using other normalizations.

From the formula (13) it follows immediately that any eigenfunction with non-zero eigenvalue is an odd function of $x - 1$ and, hence, $c(2k)(1) = 0$ for any integer $k \geq 0$, where $c(2k)(x)$ is the $2k$-th derivative of $c(x)$.

Except for $\alpha = 0$ (considered in detail below) there does not appear to be a straightforward closed-form for the eigenfunctions except in one case due to Bayly. Taking $c(x) = e^{i\alpha(x - 1)} - e^{i\alpha(x - 1)} = 2i \sin \alpha (x - 1)$, a straightforward calculation shows that (14) is satisfied for $\lambda = e^{-i\alpha}$. We refer to this example, or its normalized form $\sin(\alpha(1 - x)/\sin \alpha$, as the Bayly eigenfunction, denoted by $c_B(x)$. Although we have not proved it, the Bayly eigenfunction appears to be the only eigenfunction with eigenvalue $\lambda = e^{-i\alpha}$.

Let us now assume that $\lambda \neq 0, e^{-i\alpha}$ and let us consider Equation (14). Since $c(1) = 0$, evaluating (14) at $x = -1$ gives $e^{-i\alpha}c(-1) = e^{i\alpha}c(1) = \lambda c(-1)$, from which it follows that $c(-1) = 0$, provided $\lambda \neq e^{-i\alpha}$. Now, assuming that $3 \in D_r$, we may evaluate Equation (14) at $x = 3$, to give $e^{i\alpha}c(1) - e^{-i\alpha}c(-1) = \lambda c(3)$, so that $c(3) = 0$, although we shall not make use of this fact subsequently. Note that the Bayly eigenfunction $c_B(x)$ satisfies $c_B(1) = 0$, but $c_B(-1) \neq 0$, except for the special values $\alpha = \frac{n\pi}{2}$, $n \in \mathbb{N}$. Thus the assumption $\lambda \neq e^{-i\alpha}$ is necessary. However, all eigenfunctions $c(x)$ with eigenvalue $\lambda \neq 0, e^{-i\alpha}$ necessarily satisfy $c(1) = c(-1) = c(3) = 0$ so that a factor of $c(x)$ is $(x - 1)(x + 1)(x - 3)$.

3. Banach space formalism

Let us now turn to the Banach spaces in which we shall work. Let $r > 1$ be a real constant and let $D_r = \{ x \in \mathbb{C} : |x| < r \}$ be the disc of radius $r$ centred at 0. Let $\mathcal{F}_r$ be the complex
Banach space of complex functions $f(x)$ analytic in $D_r$ such that the $\ell_1$-norm
\[
\|f\|_1 = \sum_{k=0}^{\infty} |f_k| < \infty,
\text{ where } f(x) = \sum_{k=0}^{\infty} f_k \left( \frac{x}{r} \right)^k.
\] (15)

Where the value of $r$ is not significant, we shall write $\mathcal{F}$ and $D$ for brevity. It is possible to use other domains in $\mathbb{C}$, in particular a disc centred at a point other than 0, e.g. centred at 1, but in this paper, a disc centred at 0 is sufficient for our purposes. Note that, for consistency, we use $x$ to denote a complex variable as well as a real variable. Whichever variable is meant may be determined from context.

Although other norms are possible in principle, the $\ell_1$-norm has some properties that make it useful for numerical computation. In particular, the following hold for functions $f, g \in \mathcal{F}$ and scalars $\mu, \lambda \in \mathbb{C}$ where $f(x) = \sum_{k=0}^{\infty} f_k (x/r)^k$.

1. $\lambda f + \mu g \in \mathcal{F}$ with $\|\lambda f + \mu g\| \leq |\lambda| \|f\| + |\mu| \|g\|$.
2. $fg \in \mathcal{F}$ with $\|fg\| \leq \|f\| \|g\|$.
3. For $x \in D_r$, $|f(x)| \leq \|f\|$.
4. If $\|g\| < r$, then $f \circ g \in \mathcal{F}$ and $\|f \circ g\| \leq \|f\| (\|g\|/r)^k$ where $k = \max\{j \geq 0 : f_i = 0, i < j\}$.
5. If $\|g\|$ is less than the radius of convergence of the power series $h(x) = \sum_{k=0}^{\infty} h_k x^k$, $h_k \in \mathbb{R}$ with $h_k \geq 0$ for all $k \geq 0$, then $h \circ g \in \mathcal{F}$ and $\|h \circ g\| \leq h(\|g\|)$.
6. For $T : \mathcal{F} \to \mathcal{F}$ a bounded linear operator, the operator norm $\|T\| = \sup_{k \geq 0} \|T e_k\|$ where $e_k(x) = (x/r)^k$, the $k$th basis function. Or in other words, the operator norm $\|T\|$ is the supremum of the column norms in the matrix with respect to the basis $\{e_0, e_1, \ldots\}$.

One immediate application of these results is that the operator $S : \mathcal{F} \to \mathcal{F}$ is well defined and bounded. Indeed, $|\pm (x-1)/r| = |\pm r/2(x/r) \mp 1/2| < \|\pm r/2(x/r) \mp 1/2\| = (r + 1)/2 < r$, since $r > 1$ and so, from (12),
\[
\|Sc\| \leq 2e^{\alpha (r+1)/2} \|c\|.
\] (16)

One advantage of working in the Banach space $\mathcal{F}$ is that the operator $S$ has good spectral properties. Indeed, it follows by standard results on composition operators that, for all $\alpha \geq 0$, the operator $S = S_\alpha$ is a well-defined compact bounded linear operator $S : \mathcal{F} \to \mathcal{F}$, so that its non-zero spectrum consists of discrete eigenvalues. (See [26] for a discussion of the principles involved in the proof of these facts.)

In Sections 6 and 7 extended Banach spaces will be used, on which we place $\ell_1$-norms. In particular, we shall use the spaces $\mathbb{C} \times \mathcal{F}$ (for eigenvalue–eigenfunction pairs) and $\mathbb{C} \times \mathcal{F} \times \mathbb{C} \times \mathcal{F}$ (for eigenvalue–eigenfunction pairs and their derivatives with respect to $\alpha$, see Section 7) on which we define norms, both denoted by $\|\cdot\|$ as follows, for $(a, c) \in \mathbb{C} \times \mathcal{F}$ and $(a_1, c_1, a_2, c_2) \in \mathbb{C} \times \mathcal{F} \times \mathbb{C} \times \mathcal{F}$, by:
\[
\|(a, c)\| = |a| + \|c\|,
\] (17)
\[
\|(a_1, c_1, a_2, c_2)\| = |a_1| + \|c_1\| + |a_2| + \|c_2\|.
\] (18)

In both sections, Newton maps are defined to find solutions to eigenvalue problems.
Table 1. The eigenvalue–eigenfunction pairs for $\alpha = 0$ normalized so that the leading coefficient is 1.

<table>
<thead>
<tr>
<th>$k$</th>
<th>Degree</th>
<th>$\lambda$</th>
<th>$c(x)$</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>1</td>
<td>1/4</td>
<td>$x^1$</td>
</tr>
<tr>
<td>1</td>
<td>3</td>
<td>1/4</td>
<td>$x^5 - 5x^4 + \frac{10}{3}x^3 + 30x^2 + \frac{7}{4}x - 25$</td>
</tr>
<tr>
<td>2</td>
<td>5</td>
<td>1/64</td>
<td>$x^7 - 7x^6 - 7x^5 + 105x^4 + \frac{49}{2}x^3 - 525x^2 - \frac{31}{3}x + 427$</td>
</tr>
<tr>
<td>$k$</td>
<td>2$k$ + 1</td>
<td>1/4$k$</td>
<td>Polynomial of degree $2k + 1$</td>
</tr>
</tbody>
</table>

We now turn to a discussion of the case $\alpha = 0$, for which there is a complete theory.

4. Eigenvalue–eigenfunction pairs for $\alpha = 0$

For $\alpha = 0$, the eigenfunctions corresponding to non-zero eigenvalues are polynomials of degree $2k + 1$, for $k = 0, 1, \ldots$. Indeed we have the following theorem:

**Theorem 4.1:** On a suitably chosen Banach space $\mathcal{F}$ and for $\alpha = 0$, the operator $S$ is compact with spectrum consisting of eigenvalues $\lambda = 1/4^k$, for $k = 0, 1, 2, \ldots$ together with 0 which is itself an eigenvalue. Each positive eigenvalue $1/4^k$ is simple with eigenfunction a degree-$(2k + 1)$ polynomial $c(x)$ satisfying:

1. $c^{(2\ell)}(-1) = 0$ for $\ell = 0, 1, \ldots, k - 1$,
2. $c^{(2k)}(-1) \neq 0$,
3. $c^{(2\ell)}(1) = 0$ for $\ell = 0, 1, \ldots$.

Each even function is an eigenfunction with eigenvalue 0. The polynomial eigenfunctions for $\alpha = 0$ can be expressed in terms of Bernoulli polynomials as described below.

For $\alpha = 0$ and $k \in \mathbb{N}$ we say the eigenvalue–eigenfunction pair $(\lambda, c(x))$ is of type-$(2k + 1)$ if (1) and (2) of Theorem 4.1 are satisfied. For $\alpha > 0$, an eigenvalue–eigenfunction pair of type-$(2k + 1)$ is one contained in a continuous family of eigenvalue–eigenfunction pairs, parametrized by $\alpha$, emanating from the type-$(2k + 1)$ eigenvalue–eigenfunction pair for $\alpha = 0$.

The eigenfunctions corresponding to odd-degree polynomials for $\alpha = 0$ can be readily computed using computer algebra and the first few are given in Table 1, normalized so that each is a monic polynomial.

The eigenfunctions in Table 1 may be expressed in terms of a generating function, analogously to classical Appel sequences, the most famous of which is the sequence of Bernoulli polynomials [8]. Indeed, following the approach in [33] and [34], we may write

$$G(x, t) = \frac{2t}{e^{2t} - e^{-2t}} \left( e^{(x-1)t} - e^{-(x-1)t} \right).$$  \hspace{1cm} (19)

A standard calculation gives the identity

$$S_0 G(x, t) = G((x - 1)/2, t) - G((1 - x)/2, t) = G(x, t/2),$$ \hspace{1cm} (20)
so that, expanding $G(x, t)$ as a power series in $t$

$$G(x, t) = \sum_{n=0}^{\infty} c_n(x) \frac{t^n}{n!}, \quad (21)$$

we obtain the eigenfunctions given in Table 1 for odd $n = 2k + 1$ while $c_n(x) = 0$ for $n$ even.

The generating-function representation not only facilitates the calculation of the eigenfunctions for $\alpha = 0$ but also provides a tool to obtain many properties of these eigenfunctions, similar to the large number of properties of the Bernoulli polynomials as tabulated in [1] and online resources such as [12]. This is no coincidence; indeed there is a strong connection between the eigenfunctions of $S_0$ and the Bernoulli polynomials. Recall that the sequence of Bernoulli polynomials $B_n(x)$ is given by the generating-function expansion

$$G_B(x, t) = \frac{te^{xt}}{e^{t} - 1} = \sum_{n=0}^{\infty} B_n(x) \frac{t^n}{n!}. \quad (22)$$

Then, for odd $n$, the eigenfunctions $c_n(x)$ can be expressed in terms of $B_n(x)$:

$$c_n(x) = 2^{2n}B_n \left( \frac{x + 1}{4} \right) \quad (n \text{ odd}). \quad (23)$$

This means that many properties of the Bernoulli polynomials translate into analogous properties of the eigenfunctions $c_n(x)$ and vice versa.

From the generating function (19) some properties of $c_n(x)$ for the case $\alpha = 0$ follow straightforwardly. First each eigenfunction $c_n(x)$ is an odd function of $x−1$, as expected. Moreover, since $G(x, −t) = −G(x, t)$, the coefficient of each even power of $t$ in (21) is zero, so that $c_n(x)$ is zero for $n$ even in the expansion (21). This corresponds to the fact that every even polynomial is an eigenfunction with eigenvalue 0 for $\alpha = 0$. Moreover, for odd $n$, the function $c_n(x)$ in (21) is a polynomial of degree $n$. Indeed, for $m$ a positive integer, differentiating (19) $2m$ times gives

$$\sum_{n=0}^{\infty} c_n^{(2m)}(x) \frac{t^n}{n!} = \frac{\partial^{2m}G}{\partial x^{2m}} = t^{2m}G(x, t) = \sum_{\ell=0}^{\infty} c_\ell(x) \frac{t^{\ell+2m}}{\ell!} \quad (24)$$

$$= \sum_{n=2m}^{\infty} n(n−1) \ldots (n−2m+1) c_{n−2m}(x) \frac{t^n}{n!}, \quad (25)$$

whence $c_n^{(2m)}(x) = 0$ for $n < 2m$ and, in general, we have the relation $c_n^{(2m)}(x) = n(n−1) \ldots (n−2m+1)c_{n−2m}(x)$. Noting that $c_1(x) = x−1$, and taking $m = k$, it follows that for odd $n = 2k + 1$, $c_{2k+1}(x)$ is a monic polynomial of degree $2k + 1$.

Moreover, evaluating (19) at $x = −1$ gives $G(−1, t) = −2t$ so that $c_n(−1) = 0$ for $n \neq 1$, while $c_1(−1) = −2$ as expected. Combining this result with the previous gives immediately $c_{2k+1}^{(2m)}(−1) = 0$, for $2m < 2k$, while $c_{2k+1}^{(2k)}(−1) = −2(2k + 1)!$, $c_{2k+1}^{(2k+1)}(x) = (2k + 1)!$. This corresponds to the result of Theorem 4.1 above.
Finally, it is worth noting that for the type-\((2k + 1)\) normalized eigenfunctions \(c_{2k+1}(x)/c_{2k+1}(0) \to \cos(\pi x/2)\) as \(k \to \infty\) and there are strong results on the asymptotics of the zeros of \(c_{2k+1}(x)\) for \(\alpha = 0\). (See [11, 25] for the relevant results for the Bernoulli polynomials.)

The properties 1–3 in Theorem 4.1 are valid only for \(\alpha = 0\). But each eigenfunction corresponding to a non-zero eigenvalue is a polynomial of odd degree having precisely \(2k + 1\) roots in the complex plane, and it is possible to follow each eigenfunction as \(\alpha\) increases from 0. We refer to the family of eigenfunctions emanating from the degree-\((2k + 1)\) eigenfunction at \(\alpha = 0\) as the type-\((2k + 1)\) eigenfunctions.

In what follows our aim is to choose \(k \geq 1\) and to show that for some \(\alpha_{\text{max}}\) sufficiently large, there is a type-\((2k + 1)\) eigenvalue–eigenfunction pair for \(\alpha\) satisfying \(0 \leq \alpha \leq \alpha_{\text{max}}\), and by bounding \(|\lambda|\) to show that for \(\frac{\pi}{2} < \alpha \leq \alpha_{\text{max}}\), there exists an eigenvalue of modulus greater than 1, thereby confirming the conjecture in [17, 18].

5. Recasting the eigenvalue problem as nonlinear root-finding

As is evident, the eigenfunctions in Table 1 are monic polynomials so that the type-\((2k + 1)\) eigenfunction \(c_{2k+1}(x)\) satisfies the normalization \(c_{2k+1}^{(2k+1)}(0) = (2k + 1)!\). In general let us consider a continuous linear functional \(L : \mathcal{F} \to \mathbb{C}\) and the normalization \(L(c) = 1\). Note that the type-\((2k + 1)\) eigenfunction in Table 1 satisfies \(L_{2k+1}(c_{2k+1}) = 1\) where \(L_{2k+1}(c) = c^{(2k+1)}(0)/((2k + 1)!))\). This particular normalization is computationally expensive to implement since it involves a \((2k + 1)\)th derivative, so in practice, it is often better to use a simpler normalization such as \(L(c) = c(0)\). Note that, for a given normalization \(L\), not all functions in \(\mathcal{F}\) may necessarily be normalized since the linear functional \(L\) has a non-trivial kernel. However, for a given choice of \(k\), a suitable normalization may be identified. In the work described here, the choice \(L(c) = c(0)\) has been used.

We now assume that a normalization functional \(L\) has been specified and we define a map \(G : \mathbb{C} \times \mathcal{F} \to \mathbb{C} \times \mathcal{F}\) by

\[
G(\lambda, c) = (L(c) - 1, (S - \lambda)c).
\]

Then a normalized eigenvalue–eigenfunction pair \((\lambda, c)\) satisfies \(G(\lambda, c) = (0, 0)\), so that standard root-finding methods may be used, in particular the Newton method. Indeed, \(dG\), the derivative of \(G\) at \((\lambda, c)\) acting on tangent vectors \((\delta \lambda, \delta c) \in \mathbb{C} \times \mathcal{F}\), is given by

\[
dG_{(\lambda, c)}(\delta \lambda, \delta c) = (L(\delta c), (S - \lambda)\delta c - \delta \lambda c),
\]

so that, for fixed \(\alpha \geq 0\), the eigenvalue–eigenfunction pair can be found by the standard iteration scheme

\[
(\lambda, c) \mapsto (\lambda, c) - dG^{-1}_{(\lambda, c)} G(\lambda, c),
\]

with quadratic convergence to a simple eigenvalue–eigenfunction pair. Indeed, close to a fixed eigenvalue–eigenfunction pair, the quadratic convergence ensures that the map \(G\) is a contraction mapping. For convenience, an approximation to the Newton map, chosen so that it shares the same fixed points as the exact map and still gives a contraction close to the eigenvalue–eigenfunction pair, is used to avoid inverting the derivative \(dG\).
The Newton method may be adapted to rigorous computer-assisted functional analysis using standard interval analysis and its extension to function balls in the function space $\mathcal{F}$, as described in outline below.

6. Analysis for $\alpha \in [0, 5]$

To prove Theorem 1.1, we require two subsidiary, related theorems. The first, which is discussed in this section, gives the results on whether $|\lambda|$ is greater or less than 1 for all $\alpha \in [0, 5]$, excepting an interval around $\pi/2$, namely, $(1.5705, 1.571)$, for which less information is provided.

**Theorem 6.1 (Computer-assisted):** Let $\alpha \in [0, 5]$. Then there exists an eigenvalue–eigenfunction pair $(\lambda, c(x))$ of the operator $S$ satisfying the following:

1. $c(x) \in \mathcal{F}_r$, for some $r > 1$.
2. (a) For $\alpha \in [0, 1.5705]$, $|\lambda| < 1$.
   (b) For $\alpha \in [1.571, 5]$, $|\lambda| > 1$.
   (c) For $\alpha \in [1.5705, 1.571]$, $|\lambda| \in [0.99624, 1.00374]$.

The proof of Theorem 6.1 is based on standard rigorous computer-assisted numerical functional analysis as outlined in [13]. We describe the method of proof in the following proof algorithm. The computer estimates are described in further detail below.

6.1. **Proof algorithm**

The method of proof is to divide the interval $[0, 5]$ into a (rather large) set of subintervals such that the errors resulting by restricting $\alpha$ to each subinterval are sufficiently small for rigorous results to be obtained. In particular, a complex interval$^1$ is calculated that is guaranteed to contain an eigenvalue of the SFS operator for each $\alpha$ in the subinterval. These intervals provide the justification of the statements in Theorem 6.1. In practice, for each subinterval, high-precision calculations of an approximate eigenvalue–eigenfunction pair are performed non-rigorously for $\alpha$ set to the midpoint of the subinterval, and then rigorous calculations are performed on the whole subinterval. The parameters used in the computations are $\alpha$-dependent; in general the calculation becomes increasingly time-consuming as $\alpha$ increases and this puts a severe limitation on the maximum value of $\alpha$ for which it is possible to obtain rigorous results. A more detailed description follows.

Suppose $\alpha_0 = 0$ and for $i = 1, 2, \ldots$, we define a sequence of contiguous, possibly overlapping, real closed intervals $I_i = [\alpha_{i-1}, \alpha_i]$ such that $[0, 5] \subseteq \cup_i I_i$. For each $i$, we choose parameters, largely by trial and error, of the length $\delta \alpha_i$ of the interval $I_i$, of the radius $r_i > 1$, of the precision of floating-point number calculations and of the degree of truncation of the Taylor series expanded on the domain $D_{r_i} \supseteq [-1, 1]$. We let $\tilde{\alpha}_i \in I_i$ with $\tilde{\alpha}_i \approx (\alpha_i + \alpha_{i-1})/2$ and calculate to a high precision an approximate eigenvalue–eigenfunction pair $(\tilde{\lambda}_i, \tilde{c}_i)$, either directly or using previously calculated results for $i-1$. In practice the non-rigorous Newton method is often an efficient and convenient method to obtain such a pair. Indeed we may use (a finite-degree version of) Newton’s method (28) to obtain the
approximation to a required degree of accuracy $\delta_i$, say. This process will also provide a finite-degree approximation $\bar{d}G_{i}$ to $dG(\tilde{\lambda}_i, \tilde{c}_i)$.

In principle, many of the parameters may be allowed to vary with $i$, although in practice it is often convenient to keep the parameters largely fixed for a given run of the computer program covering a range of $\alpha$ in $[0,5]$. The choice of parameters determines the Banach spaces $\mathcal{F}_i$ and $\mathbb{C} \times \mathcal{F}_i$ as described in Section 3.

Once an accurate approximate eigenvalue–eigenfunction pair has been calculated, the next step is to use computer-assisted methods to prove rigorous bounds for the eigenvalue–eigenfunction pair $(\lambda, c)$ for all $\alpha \in I_i$. The idea is to fix $\alpha \in I_i$ and to take a (small) closed disc-function-ball $B = B_\varepsilon(\tilde{\lambda}_i, \tilde{c}_i)$ around $(\tilde{\lambda}_i, \tilde{c}_i)$ in $\mathbb{C} \times \mathcal{F}_i$ (equipped with a norm $\|\|$) and use an approximation to the Newton map to define a contraction on $B$ using the following idea. Let $N_i$ be the approximate Newton map

$$N_i : (\lambda, c) \mapsto (\lambda, c) - F_i G(\lambda, c)$$

where $F_i$ is a fixed linear operator which approximates the inverse of $dG(\tilde{\lambda}_i, \tilde{c}_i)$ evaluated at $\tilde{\alpha}_i$. Note that $N_i$ does not have the quadratic convergence of the Newton method, but nevertheless should be a contraction, provided the parameters are chosen suitably. In fact, the derivative $dN_i(\lambda, c) = I - F_i dG(\lambda, c)$ should be close to zero and certainly satisfy $\sup_{(\lambda, c) \in B} \|dN_i(\lambda, c)\| \leq \eta < 1$, where $\|dN_i(\lambda, c)\|$ is the induced operator norm.

Now suppose $\|N_i(\tilde{\lambda}_i, \tilde{c}_i) - (\tilde{\lambda}_i, \tilde{c}_i)\| < \delta$. Then, for $(\lambda, c) \in B$,

$$\|N_i(\lambda, c) - (\tilde{\lambda}_i, \tilde{c}_i)\| \leq \|N_i(\tilde{\lambda}_i, \tilde{c}_i) - (\tilde{\lambda}_i, \tilde{c}_i)\| + \|N_i(\lambda, c) - N_i(\tilde{\lambda}_i, \tilde{c}_i)\|$$

$$\leq \|N_i(\tilde{\lambda}_i, \tilde{c}_i) - (\tilde{\lambda}_i, \tilde{c}_i)\| + \sup_{(\lambda, c) \in B} \|dN_i(\lambda, c)\| \|\|(\lambda, c) - (\tilde{\lambda}_i, \tilde{c}_i)\|$$

$$\leq \delta + \eta \varepsilon \leq \varepsilon,$$

provided $\delta \leq (1 - \eta)\varepsilon$, a condition that must be checked. This means that $N_i$ maps $B \rightarrow B$ and is clearly a contraction (since $\eta < 1$) and so has a unique fixed point, for a fixed $\alpha \in I_i$. This fixed point gives the eigenvalue–eigenfunction pair. That $N_i$ is a contraction, and therefore each eigenvalue–eigenfunction pair is unique within the ball $B_i$, will be an important part of the argument for the results in the next section.

In this way, using several independent runs of the computer program, it is possible to obtain results for all the $I_i$ and thereby to cover the whole of $[0,5]$. This provides the proof of Theorem 6.1.

7. Analysis around $\alpha = \pi/2$

The computer-assisted results described in the previous sections do not cover in sufficient detail the values of $\alpha$ close to $\pi/2$ and, in this section, we describe additional results which cover an interval of $\alpha$ straddling $\alpha = \pi/2$. Specifically, we describe the proof of the following theorem:
**Theorem 7.1 (Computer-assisted):** Let \( \alpha \in [1.5705, 1.571] \). Then there exists an eigenvalue–eigenfunction pair \((\lambda, c(x))\) of the operator \( S \) satisfying the following:

1. \( \lambda \in [-0.00375, 0.00375] + i[-1.00375, -0.99625] \) with \( |\lambda| \in [0.99624, 1.00374] \)
2. \( c(x) \in \mathcal{F}_{1.1}, c(1) = c(-1) = 0, c(0) = 1. \)
3. (a) For \( \alpha < \pi/2, |\lambda| < 1. \)
   (b) For \( \alpha > \pi/2, |\lambda| > 1. \)
   (c) For \( \alpha = \pi/2, \lambda = -i. \)

Combining this theorem with Theorem 6.1 gives Theorem 1.1. Again, we note that 3(a) makes no claim about all eigenvalues, only about a single eigenvalue found through the computer-assisted proof.

The proof of Theorem 7.1 requires a number of steps as described below. In particular, because of the existence of the Bayly family \( c_B(x) \) of eigenfunctions described in Section 2, which crosses the type-3 family of eigenfunctions at \( \alpha = \pi/2 \), the function \( c(x) \) is required to have zeros at \( x = -1, 1 \) to distinguish it from the Bayly family, so that the representation \( c(x) = (x + 1)(x - 1)\hat{c}(x) \) has been used to ensure these conditions are satisfied. We describe the approach to proving Theorem 7.1 in a series of steps. The essential idea is to prove that, for each \( \alpha \) in the interval \([1.5705, 1.571]\), there exists a unique eigenvalue–eigenfunction pair in a closed ball in the Banach space \( \mathbb{C} \times \mathcal{F}_{1.1} \) (\( r \) is set to 1.1 for this work) and to obtain bounds on \( d\lambda/d\alpha \) for \( \alpha \) in this interval, to show that \( d|\lambda|/d\alpha > 0 \), while at the same time showing that the Bayly eigenfunction for \( \alpha = \pi/2 \) (suitably normalized) lies in this closed ball. This latter fact shows that the Bayly eigenfunction is in fact the unique eigenfunction in the ball for \( \alpha = \pi/2 \) and since it is already established that \( \lambda = -i \) for the Bayly eigenfunction, the precise statement on \( \lambda \) in Theorem 7.1 is proved.

Let us write \( c(x) = (x - 1)(1 + x)\hat{c}(x) \), where \( \hat{c}(x) \in \mathcal{F}_r \), for \( r > 1 \). Then the operator \( S \) induces an operator \( \hat{S} \) on \( \hat{c}(x) \) as follows

\[
S((x - 1)(1 + x)\hat{c}(x)) = e^{i\alpha(x-1)2} \left( \frac{x - 1}{2} - 1 \right) (x - 1) + 1)\hat{c}\left(\frac{x - 1}{2}\right) - e^{i\alpha\left(\frac{x}{2}-\right)2} \left( \frac{1 - x}{2} - 1 \right) (1 - x) + 1\hat{c}\left(\frac{1 - x}{2}\right) = e^{i\alpha(x-1)2} \left( \frac{x - 3}{2} \right) (x + 1)\frac{2}{2})\hat{c}\left(\frac{x - 1}{2}\right) - e^{i\alpha\left(\frac{x}{2}-\right)2} \left( \frac{-1 - x}{2} \right) (3 - x)\frac{2}{2})\hat{c}\left(\frac{1 - x}{2}\right) = \frac{(x + 1)(x - 1)(x - 3)}{4(x - 1)} \left[ e^{i\alpha(x-1)2} \hat{c}\left(\frac{x - 1}{2}\right) - e^{i\alpha\left(\frac{x}{2}-\right)2} \hat{c}\left(\frac{1 - x}{2}\right) \right]
\]

which implies

\[
\hat{S}\hat{c}(x) = \frac{(x - 3)}{4(x - 1)} S\hat{c}(x) = \frac{(x - 3)}{8(x - 1)/2} \left[ e^{i\alpha x} \hat{c}(x) \right]_{2,0} \left( \frac{x - 1}{2} \right).
\]
Note that the term $x-1$ in the denominator of Equation (30) cancels with $x-1$ in the term $\hat{S}\hat{c}(x)$ because the latter is an odd function of $x-1$. Modulo any function-space technicalities, the operators $\hat{S}$ and $\hat{S}$ have the same spectrum when $\hat{S}$ is restricted to functions $c(x)$ satisfying $c(\pm1) = 0$.

Now let us consider the operator $\hat{S}$ given by Equation (30) on $\mathcal{F}_r$. Noting that for $\hat{\chi}(x) \in \mathcal{F}_r$, $e^{i\alpha x}\hat{\chi}(x) \in \mathcal{F}_r$ with norm bounded by $e^{\alpha r}||\hat{\chi}||$, and $[e^{i\alpha x}\hat{\chi}(x)]_2/\chi \in \mathcal{F}_r$ with norm bounded by $2e^{\alpha r}||\hat{\chi}||$, it follows that $\hat{S}\hat{\chi}(x) \in \mathcal{F}_r$ with norm

$$||\hat{S}\hat{\chi}(x)|| \leq \frac{3+r}{4} e^{\alpha r}||\hat{\chi}||.$$  

(31)

We note that if $\hat{\chi}(x) \in \mathcal{F}_r$ then $c(x) = (x-1)(x+1)\hat{\chi}(x) \in \mathcal{F}_r$ so that an eigenfunction $\hat{\chi}(x)$ of $\hat{S}$ in $\mathcal{F}_r$ gives an eigenfunction $c(x)$ of $S$ also in $\mathcal{F}_r$.

For general $\alpha$ the Bayly eigenfunction $c_B(x)$ cannot be written in this way except for certain values of $\alpha$ (specifically $\alpha = n\pi/2$, $n \in \mathbb{N}$). Indeed, for $\alpha = \pi/2$, using the Euler product formula for the sine function

$$\sin \pi x = \pi x \prod_{k=1}^\infty \left(1 - \frac{x^2}{k^2}\right)$$  

(32)

gives

$$\sin \frac{(x-1)}{2} = \frac{\pi (x-1)}{2} \prod_{k=1}^\infty \left(1 - \frac{(x-1)^2}{4k^2}\right) = (x-1)(x+1)\hat{\chi}(x),$$  

(33)

where

$$\hat{\chi}(x) = \frac{\pi (3-x)}{8} \prod_{k=2}^\infty \left(1 - \frac{(x-1)^2}{4k^2}\right).$$  

(34)

The steps in detail are enumerated below.

1. An accurate approximate eigenvalue–eigenvalue pair $(\lambda_0, c_0(x)) \in \mathbb{C} \times \mathcal{F}_{1,1}$ is calculated for $\alpha$ close to $\pi/2$ and a ball $B$ around $(\lambda_0, c_0(x))$ is defined in $\mathbb{C} \times \mathcal{F}_{1,1}$.

2. For each $\alpha \in [1.5705, 1.571]$, it is shown, as in Section 6, that an approximate Newton map $N$ is a contraction on $B$ so that there is a unique eigenvalue–eigenvector pair $(\lambda_*, c_*(x))$ (dependent on $\alpha$).

3. Using the expansion (34), it is shown that the Bayly eigenfunction pair corresponding to $(-i, \sin \pi (1-x)/2)$ lies in $B$. This is a non-trivial computation for which the expansion (33) is used. Indeed, by truncating (34) at a suitably high, but finite, number of terms in the infinite product, and estimating the error in so doing, it is possible to show that the Bayly eigenfunction lies in a function ball contained in $B$. Because of the uniqueness of $(\lambda_*, c_*(x)) \in B$ for each $\alpha$, it follows that, for $\alpha = \pi/2$, the pair $(\lambda_*, c_*(x))$ corresponds to $(-i, \sin \pi (1-x)/2)$.

4. From the implicit function theorem, it follows that $(\lambda_*, c_*)$ depends smoothly on $\alpha$. Writing $\hat{\lambda} = \partial \lambda_*/\partial \alpha$ and $\hat{\chi}(x) = \partial c_*(x)/\partial \alpha$, the next step is to find a ball $V$ containing the pair $(\hat{\lambda}, \hat{\chi}(x))$. The method used is an adaptation of the approximate-Newton method used to obtain the eigenvalue–eigenvector pair $(\lambda_*, c_*(x))$. 
Indeed, we define a map $H$ on a ball in $\mathbb{C} \times \mathcal{F}_{1.1} \times \mathbb{C} \times \mathcal{F}_{1.1}$ to itself by

$$H(\lambda, c, \dot{\lambda}, \dot{c}) = (L(c) - 1, (S - \lambda)c, L(\dot{c}), \dot{S}(c) + S(\dot{c}) - \dot{\lambda}c - \lambda\dot{c})$$

(35)

where $\dot{S}$ is the operator $\partial S/\partial \alpha$, and seek a zero of $H$ using an approximate Newton method as in Section 6.

(5) Taking the bounds on values of $\dot{\lambda}$ from $\mathcal{V}$ and on $\lambda$ from $\mathcal{B}$, the value of $d|\lambda|^2/d\alpha$ is estimated for all $\alpha \in [1.5705, 1.571]$. Calculating bounds on $d|\lambda|^2/d\alpha$ is straightforward; using the relation $d|\lambda|^2/d\alpha = (d\lambda/d\alpha)\dot{\lambda} + \lambda(d\dot{\lambda}/d\alpha)$ we find that $d|\lambda|^2/d\alpha$ is contained in a strictly positive real interval $[0.7466, 0.7673]$. It follows that $|\lambda|$ increases as a function of $\alpha$ in this range and since $|\lambda| = 1$ for $\alpha = \pi/2$, it follows that, for $\alpha \in [1.5705, 1.571], |\lambda| < 1$ for $\alpha < \pi/2$ and $|\lambda| > 1$ for $\alpha > \pi/2$.

8. Computer-assisted estimates

The results in this paper depend on computer-assisted analytic estimates and, in this section, we outline the methodology we have used to obtain these estimates.

The starting point for computer-assisted analytic estimates is the classical theory of interval arithmetic [2, 20, 27–29] with rigorous error bounding to take into account computer rounding error. For finite closed intervals of real numbers, one defines unary and binary operations corresponding to the usual algebraic operations with output a finite closed interval of real numbers which contains the set of all results of performing the operations pointwise on the intervals, in such a way that the output interval is close to optimal. The Julia programming language [7] with the interval-arithmetic package IntervalArithmetic.jl [6] is a convenient programming environment that implements fixed (but arbitrary) precision interval arithmetic with an extension to intervals of complex numbers [24, 29].

To obtain the functional analytic estimates, we use an extension of interval arithmetic to balls in function space, an approach that has been used extensively over the last 40 years [9, 13, 14, 22, 23, 32]. Analytic functions are completely determined by their Taylor series about the centre of their domain. For computations in a computer with a finite amount of storage space, the Taylor series can be truncated after a certain number of terms. Detailed information on the Taylor polynomial of degree $N$ (say) can be retained (by specifying complex intervals in which the Taylor coefficients lie); however, we must retain information on the norm of any higher-order terms. In fact, because certain operations on higher-order terms result in lower-order terms it makes sense to include an additional 'error' function to account for such terms, for which a bound on the may be kept. These considerations lead to the following definition of a complex function ball in $\mathcal{F}_r$ [13].

Fixing $r > 1$ and the function space $\mathcal{F}_r$, for a given (fixed) positive integer $N$ (the degree of truncation) and a sequence $I_{P,i}, i = 0, \ldots, N$ of $N + 1$ intervals of complex numbers and two non-negative real numbers $V_H$ and $V_E$, we may define a function ball $B \subseteq \mathcal{F}_r$ to be the set of all functions $f \in \mathcal{F}_r$ which can be written in the form

$$f(x) = f_P(x) + f_H(x) + f_E(x)$$

(36)
where $f_P(x)$, the polynomial part of lower order, is a function

$$
    f_P(x) = \sum_{i=0}^{N} f_{P,i} \left(\frac{x}{r}\right)^i, \quad f_{P,i} \in I_{P,i}, \quad i = 0, 1, \ldots, N,
$$

$f_H(x)$, the higher-order part, is a function such that

$$
    f_H(x) = \sum_{i=N+1}^{\infty} f_{H,i} \left(\frac{x}{r}\right)^i \quad \text{and} \quad \|f_H\|_1 = \sum_{i=N+1}^{\infty} |f_{H,i}| \leq V_H,
$$

and $f_E(x)$, the error part, is a function such that

$$
    f_E(x) = \sum_{i=0}^{\infty} f_{E,i} \left(\frac{x}{r}\right)^i \quad \text{and} \quad \|f_E\|_1 = \sum_{i=0}^{\infty} |f_{E,i}| \leq V_E.
$$

In this definition, we have detailed information about $f_P$ since its Taylor coefficients lie in complex intervals that are known explicitly. But $f_H$ and $f_E$ are only known roughly by the bounds on their $\ell_1$-norms. The error function $f_E$ is introduced so that when lower-order terms are produced from higher-order terms as a result of, say, differentiating a higher-order function, we will not have detailed information about the coefficients, so that we may allocate the result to the error function, even if the terms are indeed of low order.

The ideas of interval arithmetic may be extended to function balls, so that binary and unary operations on functions in $F_r$ may be extended to operations on function balls, and implemented on a computer, by adapting the operations to take into account computer rounding error. To achieve this, a suite of Julia programs was written, based on the IntervalArithmetic.jl package, to implement the algorithms described in Sections 6.1 and 7.

For the code that implements the computer proof see [30] and [31]. In practice, we have found that taking $r = 1.1$, polynomials of degree $N$ between 10 and 25, and Float64-precision has been sufficient to prove the results given in this article. The required choice of interval length for the parameter $\alpha$ varies with the size of $\alpha$. For the proofs contained herein, interval lengths varying between $10^{-5}$ (for small values of $\alpha$) and $10^{-4}$ (for $\alpha$ around 5.0) have been sufficient. Optimisation of the algorithm and code would likely increase the interval lengths.

9. Conclusion and further work

The work described in this article constitutes a preliminary study of the Stretch-Fold-Shear (SFS) operator for $\alpha \in [0, 5]$. For $\alpha = 0$, the spectrum is readily determined and the polynomial eigenfunctions may be obtained using generating-function techniques; the eigenfunctions are strongly connected with the classical Bernoulli polynomials. Despite the apparent simplicity of the SFS operator, the determination of its spectrum for $\alpha > 0$ is a non-trivial endeavour and, although the spectrum has been approximated numerically, the computer-assisted results presented in this article are the first rigorous results to be obtained.

In this article, we have described a computer-assisted proof to obtain rigorous bounds for $\alpha \in [0, 5]$ on the eigenvalue of the type-3 eigenfunction, which corresponds, for $\alpha = 0,$
to a degree-3 polynomial eigenfunction. The methods can be readily adapted to other odd
types (5, 7, etc.), and for larger values of $\alpha$.

A word of caution is required here. The computer-assisted methods require accurate
computations that become increasingly time consuming for higher eigenfunction types
and even for relatively small values of $\alpha$. Hence, while powerful in themselves, they do not
appear to be applicable to the study of the spectrum for large values of $\alpha$, let alone the
determination of the asymptotic properties of the spectrum as $\alpha \to \infty$.

The generating-function approach is an elegant technique that provides powerful results
for $\alpha = 0$. It is unclear to what extent these techniques can be used for $\alpha > 0$ and the
investigation of their generalization to $\alpha > 0$ is an important area of future research in this
field.

**Note**

1. Here, and in what follows unless explicitly stated otherwise, a complex interval is a rectangle in
the complex plane of the form $[a, b] + i[c, d]$, where $[a, b]$ and $[c, d]$ are real intervals. Interval
arithmetic can readily be extended to complex intervals.

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**Data availability statement**

No data were created or analysed in this study. Details of the Julia programs used in the
proofs reported in this paper are available from [30] and from [31].

**References**

[1] Abramowitz M. and Stegun I.A., Handbook of Mathematical Functions with Formulas, Graphs,
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